

Dijkstra Alg for Single-Source Shortest Path in Graphs (Dir or unDir)

Given: Graph (V, E) directed or undirected
Source vertex S | assume connected

Each edge has positive weight

$$w(u, v) \geq 0$$

Goal: Produce the Dijkstra Tree (root s)

to all reachable vertices in the graph

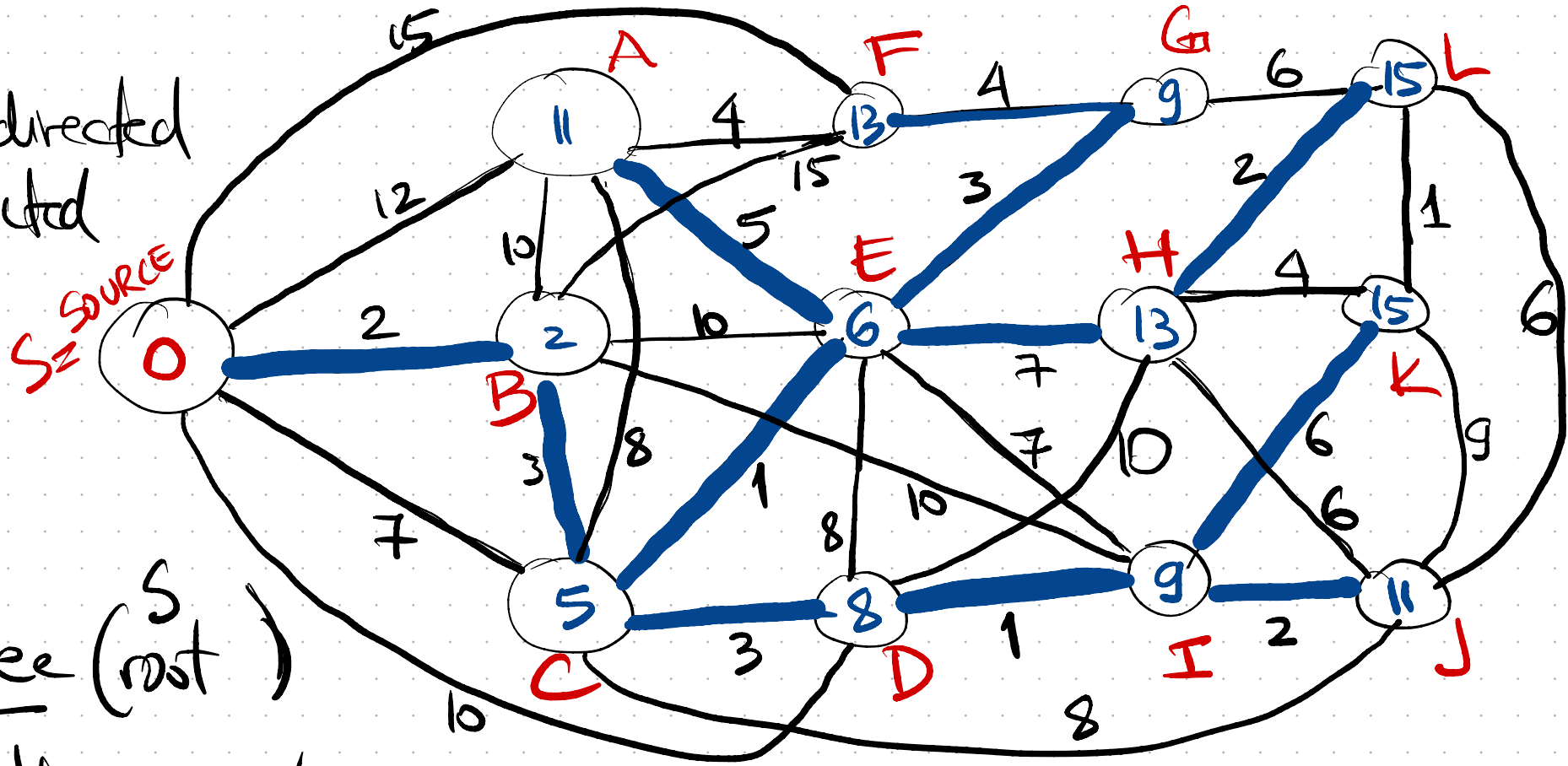
such that the tree-path to every vertex

a SP (s, u) in the graph.

Shortest path

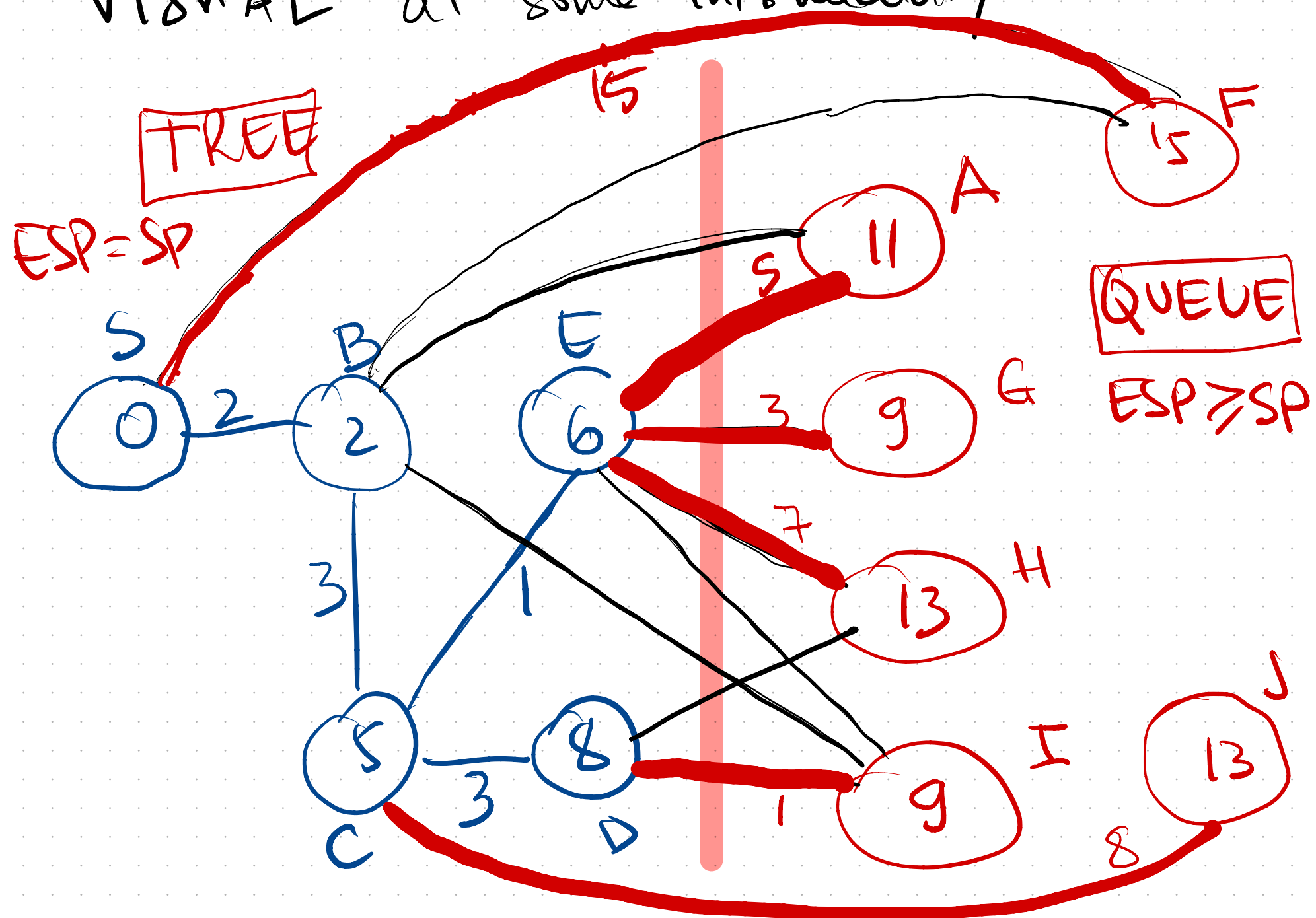
Dijkstra Algorithm: at all times vertices are partitioned into 3 groups:

TREE	QUEUE	NOT YET
already (SP done) in Dij-tree	candidates (with ESP) to add to Dij-tree	not connected yet to tree $ESP = \infty$



At any time $SP(s, u) = \text{true SP in the graph (unknown)}$
 written value $\leftarrow$ $ESP(s, u) = \text{the SP found so far. } \boxed{ESP(s, u) \geq SP(s, u)}$
 inside each node
 TREE vertices: $SP(s, u) = ESP(s, u)$

VISUAL at some intermediary time in the Algorithm:



NOT YET $ESP = \infty$

Dijkstra Theorem : • The min-ESP in the queue is actually SP
(in the example G and I have min ESP=9)

- Then the node in queue with min ESP is added to Tree with parent/edge that give the SP-ESP. (in example G can be added with parent E (edge EG) or I can be added with parent D (edge DI))

Algorithm: • maintain the queue, the Tree and update all ESPs.

- When min-ESP in queue is added (ex edge $E \rightarrow G$)
 - $G.SP = G.ESP$; $G.parent = E$
 - For all $u \in G.adj-list$: update $u.ESP$ if necessary, if $u.ESP > G.SP + w(G \rightarrow u)$ then $u.ESP = G.SP + w(G \rightarrow u)$; $u.parent = G$

Proof: Assume (hypothetical) that $\min \text{ESP}$ in the queue is not SP node u .

TREE

QUEUE

Let v be the node in the tree that connects u for $\min \text{ESP}$.
 so $v.SP + w(v \rightarrow u) = u.ESP$.

Since $u.ESP > u.SP$ let's look at the path $S \rightarrow u$ that give the SP = shortest path:

$S \rightarrow k_0 \rightarrow k_1 \rightarrow \dots \rightarrow k_n \rightarrow u$

then $u.SP = w(S \rightarrow k_0) + w(k_0 \rightarrow k_1) + \dots + w(k_{n-1} \rightarrow k_n) + w(k_n \rightarrow u)$

key: at some point this path crosses from TREE group to QUEUE
 In example that's $k_2 \rightarrow k_3$. Then look at $k_3.ESP$

$k_3.ESP < u.SP < u.ESP$

CONTRADICTION: $u.ESP$ is not the $\min\text{-ESP}$ in QUEUE.

