

Last time

- Expectation

Today

- Applications of Expectation
 - Variance & Standard Deviation
 - Entropy

Next time

- Conditional Prob.
- Bayes Theorem

Variance

Consider heights

case 1

4'10" 5' 5'2"

$$E\{x\} = 5'$$

case 2

4' 5' 6'

$$E\{x\} = 5'$$

case 3

3' 5' 7'

$$E\{x\} = 5'$$

How to measure "variability" 3 ways

case 2

$$\textcircled{1} E\{Y\} = \sum_{w \in \Omega} Y(w) \cdot p(w)$$

$$= (-12'') \cdot \frac{1}{3} + (0'') \cdot \frac{1}{3} + (+12'') \cdot \frac{1}{3}$$
$$= 0$$

$$\textcircled{2} E\{Y\} = |-12''| \cdot \frac{1}{3} + |0''| \cdot \frac{1}{3} + |12''| \cdot \frac{1}{3}$$
$$= 12 \cdot \frac{1}{3} + 0 \cdot \frac{1}{3} + 12 \cdot \frac{1}{3} = 8 \text{ inches}$$

$$\textcircled{3} E\{Y\} = E\{(X - E\{X\})^2\} = \text{case 1}$$
$$= (-2'')^2 \cdot \frac{1}{3} + (0'')^2 \cdot \frac{1}{3} + (+2'')^2 \cdot \frac{1}{3}$$
$$= \frac{4}{3} + 0 + \frac{4}{3} = \frac{8}{3} \text{ inches}^2$$

X neg. & pos. cancel

- mean absolute deviation

- variance

variance σ^2

Take square root, get back inches...

$$\sigma^2 = \text{Var}(X) = E[(X - E\{X\})^2]$$

$$\Rightarrow \sigma = \sqrt{\text{Var}(X)} = \sqrt{E[(X - E\{X\})^2]} \quad - \text{ back in units originally measured.}$$

e.g. case 1: $\sigma^2 = \text{Var}(X) = 8/3 \text{ inches}^2$

$$\sigma = \sqrt{\text{Var}(X)} = \sqrt{8/3 \text{ inches}^2} = 1.63''$$

case 3: ... $\sigma = 19.6''$

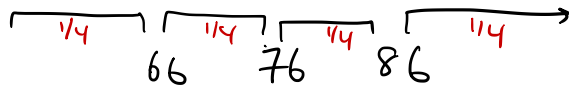
Exam

$$E\{X\} = 76$$

$$\sigma = 15$$

$\pm 2/3 \sigma$ - quantiles

$$2/3 \sigma = 10 \text{ pts}$$



Entropy

Consider 8 letters only $\{A, B, C, D, E, F, G, H\}$

Need 3-bits

A	→	000
B	→	001
C	→	010
D	→	011
E	→	100
F	→	101
G	→	110
H	→	111

} fixed length code

suppose must encode n -things, need: $\lceil \log_2 n \rceil$

Efficiency of code is measured in bits-per-character on average, BPC. $BPC = 3$

variable length code: assign short codes to frequent letters
long codes to infrequent letters.

$$BPC = \sum_{w \in \Sigma} X(w) \cdot P(w) = \sum_i l_i \cdot p_i$$

Why must we have long codes to make up for short codes?

Analog to PHP: Kraft's Inequality

Example:

A	B	C	D	E	F	G	H	BBB	→	010101
00	01	010	011	100	101	110	111	CF	→	010101

code length for i th letter

probability of i th letter

Suppose we have following skewed distribution

$$P_i = \left(\overset{A}{1/4}, \overset{B}{1/4}, \overset{C}{1/8}, \overset{D}{1/8}, \overset{E}{1/16}, \overset{F}{1/16}, \overset{G}{1/16}, \overset{H}{1/16} \right)$$

$$(00, 01, 100, 101, 1100, 1101, 1110, 1111)$$

$$ADF \rightarrow 001011101$$

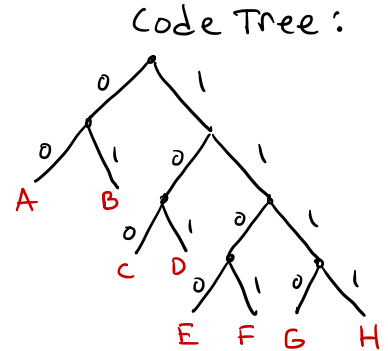
$$\begin{array}{c} 00|10|11101 \\ A \quad D \quad F \end{array}$$

$$BPC = \sum_i l_i \cdot p_i$$

$$= 2 \cdot 1/4 + 2 \cdot 1/4 + 3 \cdot 1/8 + 3 \cdot 1/8 + 4 \cdot 1/16 + \dots + 4 \cdot 1/16$$

$$= 2.75$$

* 8.33% savings over fixed code of length 3.



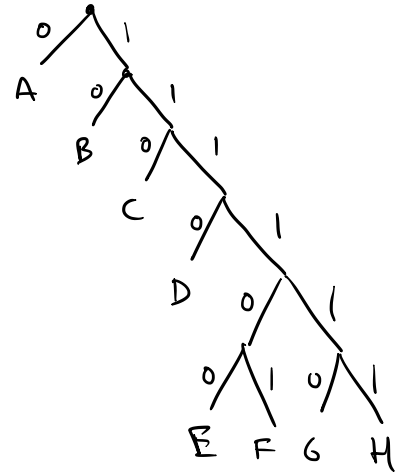
$$P_2 = \begin{matrix} A & B & C & D & E & F & G & H \\ \left(\frac{1}{2}, \frac{1}{4}, \frac{1}{8}, \frac{1}{16}, \frac{1}{64}, \frac{1}{64}, \frac{1}{64}, \frac{1}{64} \right) \end{matrix}$$

0 10 110 1110 111100 111101 111110 111111

$$BPC = \sum_i l_i \cdot p_i$$

$$= 1 \cdot \frac{1}{2} + 2 \cdot \frac{1}{4} + 3 \cdot \frac{1}{8} + 4 \cdot \frac{1}{16} + 6 \cdot \frac{1}{64} + \dots + 6 \cdot \frac{1}{64}$$

$$= 2$$



p_i	l_i
$\frac{1}{2}$	1
$\frac{1}{4}$	2
$\frac{1}{8}$	3
$\frac{1}{16}$	4
$\frac{1}{64}$	6

$$p_i = \frac{1}{2^k} \Rightarrow p_i = \frac{1}{2^{l_i}}$$

$$\Leftrightarrow 2^{l_i} = \frac{1}{p_i}$$

$$\Leftrightarrow l_i = \log_2 \frac{1}{p_i}$$

Entropy.

$$Q: \quad P = (p_1, p_2, p_3, \dots)$$

$$l_i = \log_2 \frac{1}{p_i}$$

What is expected code length?

$$BPC = \sum_i p_i \cdot l_i$$

$$= \sum_i p_i \cdot \log_2 \frac{1}{p_i}$$

Entropy

$$H(P) = \sum_i p_i \cdot \log_2 \frac{1}{p_i}$$

ACBADAFGAH

$$P_A = \frac{4}{10}$$

$$P_B = \frac{1}{10}$$

$$P_C = \frac{1}{10}$$

$$P_D = \frac{1}{10}$$

⋮

Facts

- ① Entropy is the compression limit.
- ② Entropy is essential to bound an ability to communicate in presence of noise.
- ③ Fundamental measure of randomness.
- ④ Cryptography

Conditional Probability

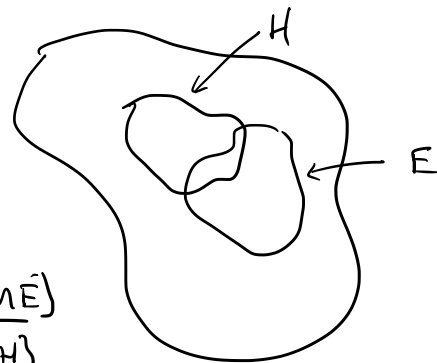
← prob of E
given H

$$Pr[H \cap E] = Pr[H] \cdot \underline{Pr[E|H]} \quad \leftarrow$$

$$\Rightarrow Pr[E|H] = \frac{Pr[H \cap E]}{Pr[H]}$$

$$= Pr[E] \cdot \underline{Pr[H|E]} \quad \leftarrow$$

$$\Rightarrow Pr[H|E] = \frac{Pr[H \cap E]}{Pr[E]}$$



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Bayes Law:

$$\underline{Pr[E]} \cdot \underline{Pr[H|E]} = Pr[H \cap E] = \underline{Pr[H]} \cdot \underline{Pr[E|H]}$$

$$Pr[H|E] = \frac{Pr[E|H] \cdot Pr[H]}{Pr[E]}$$

H = hypothesis: do have  
Zika?

E = evidence: did blood  
test come  
back pos?

CS1800  
Discrete Structures  
Fall 2017

Lecture 25  
11/2/17

## Last time

- Finish Expectation
  - Variance / standard deviation
- Start Entropy

## Today

- Finish Entropy
- Conditional Prob. & Bayes Law

## Next time

- Markov chains & Page Rank

## Announcements

Regrades → Kevin Gold

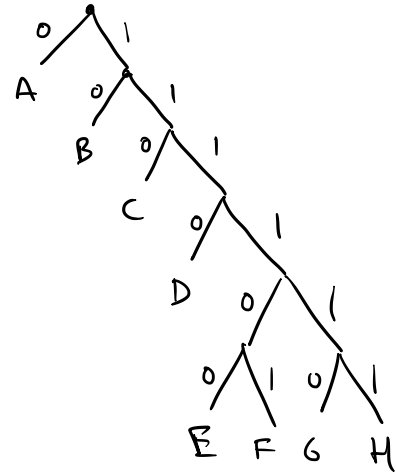
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$$BPC = \sum_i l_i \cdot p_i$$

$$= 1 \cdot \frac{1}{2} + 2 \cdot \frac{1}{4} + 3 \cdot \frac{1}{8} + 4 \cdot \frac{1}{16} + 6 \cdot \frac{1}{64} + \dots + 6 \cdot \frac{1}{64}$$

$$= 2$$



| $p_i$          | $l_i$ |
|----------------|-------|
| $\frac{1}{2}$  | 1     |
| $\frac{1}{4}$  | 2     |
| $\frac{1}{8}$  | 3     |
| $\frac{1}{16}$ | 4     |
| $\frac{1}{64}$ | 6     |

$$p_i = \frac{1}{2^k} \Rightarrow p_i = \frac{1}{2^{l_i}}$$

$$\Leftrightarrow 2^{l_i} = \frac{1}{p_i}$$

$$\Leftrightarrow l_i = \log_2 \frac{1}{p_i}$$

## Entropy.

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$$l_i = \log_2 \frac{1}{p_i}$$

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Entropy

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⋮

### Facts

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## Conditional Probability

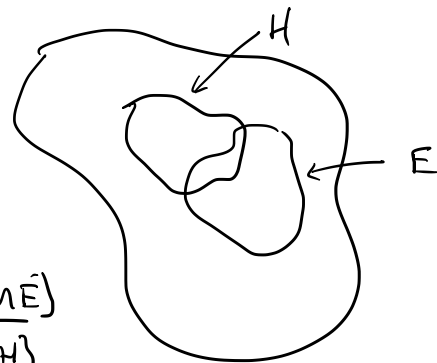
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Bayes Law:

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H = hypothesis: do have
Zika?

E = evidence: did blood
test come
back pos?

CS1800
Discrete Structures
Fall 2017

Lecture 26
11/6/17

Last time

- Finish Entropy
- Start Cond. Prob.
 - Bayes Law

Today

- Finish Cond. Prob.
 - & Bayes Law
- Start Markov chains
 - & Page Rank

Next time

- Finish MC & PR

-
- Start Module 4:
Algorithms &
Analysis

Conditional Probability

← prob of E
given H

$$Pr[H \cap E] = Pr[H] \cdot \underline{Pr[E|H]} \quad \leftarrow$$

$$\Rightarrow Pr[E|H] = \frac{Pr[H \cap E]}{Pr[H]}$$

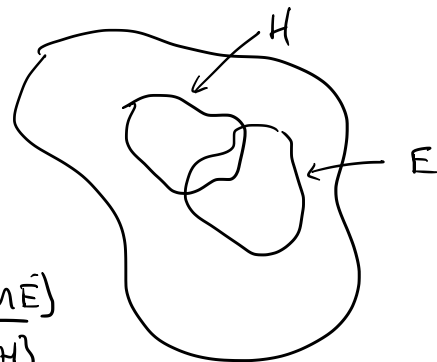
$$= Pr[E] \cdot Pr[H|E] \quad \leftarrow$$

$$\Rightarrow Pr[H|E] = \frac{Pr[H \cap E]}{Pr[E]}$$

~~~~~  
Bayes Law:

$$\underline{Pr[E]} \cdot \underline{Pr[H|E]} = Pr[H \cap E] = \underline{Pr[H]} \cdot \underline{Pr[E|H]}$$

$$Pr[H|E] = \frac{Pr[E|H] \cdot Pr[H]}{Pr[E]}$$



Independence

$$Pr[E|H] = Pr[E]$$

$$Pr[H|E] = Pr[H]$$

$$\Rightarrow Pr[H \cap E] =$$

$$Pr[H] \cdot Pr[E|H]$$

$$= Pr[H] \cdot Pr[E]$$

H = hypothesis: do have Zika?

E = evidence: did blood test come back pos?

## Bayes Law

$$Pr(B|A) \cdot Pr(A) = Pr(A \cap B) = Pr(A|B) \cdot Pr(B)$$

$$Pr(A|B) = \frac{Pr(B|A) \cdot Pr(A)}{Pr(B)}$$

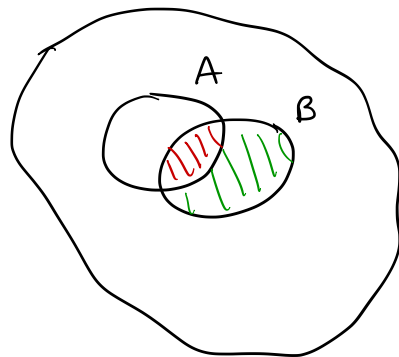
$$= \frac{Pr(B|A) \cdot Pr(A)}{Pr(B|A) \cdot Pr(A) + Pr(B|\bar{A}) \cdot Pr(\bar{A})}$$

$$Pr(B|A) \cdot Pr(A) + Pr(B|\bar{A}) \cdot Pr(\bar{A})$$

what is  $Pr(B)$ ?

$$P(B) = \text{red} + \text{green}$$

$$= P(B|A) \cdot P(A) + P(B|\bar{A}) \cdot P(\bar{A})$$



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$$Pr(H|E) = \frac{Pr(E|H) \cdot Pr(H)}{Pr(E|H) \cdot Pr(H) + Pr(E|\bar{H}) \cdot Pr(\bar{H})}$$

$$Pr(E|H) \cdot Pr(H) + Pr(E|\bar{H}) \cdot Pr(\bar{H})$$

Example: Zika in FL 2016

• Prevalence of Zika in S. FL.  $p(\text{Zika}) = 10^{-5}$  (1 in 100,000)

• accuracy of blood test is 99%

i.e.  $p(\text{pos. test} | \text{Zika}) = 0.99$  ← test subjects who had Zika

$p(\text{pos. test} | \text{no Zika}) = 0.01$  ← control group who don't have Zika

• You test positive : what is the chance that you have Zika?

⇒ Not  $p(\text{pos. test} | \text{Zika}) = 0.99$

⇒ you want  $p(\text{Zika} | \text{pos. test})$  !

$$\begin{aligned}
 P(\text{zika} \mid \text{pos test}) &= \frac{P(\text{pos. test} \mid \text{zika}) \cdot P(\text{zika})}{P(\text{pos test})} \\
 &= \frac{P(\text{pos. test} \mid \text{zika}) \cdot P(\text{zika})}{P(\text{pos test} \mid \text{zika}) \cdot P(\text{zika}) + P(\text{pos test} \mid \text{not zika}) \cdot P(\text{not zika})}
 \end{aligned}$$

$$= \frac{0.99 \times 10^{-5}}{0.99 \times 10^{-5} + 0.01 \times (1 - 10^{-5})}$$

$$= \frac{0.0000099}{0.0000099 + 0.0099999}$$

$$\approx 0.00099$$

$$\approx 0.1\% \quad \text{i.e., only 1 in 1,000!}$$

Seems wildly counter intuitive, but...

10,000,000 people in FL.  $10^7$

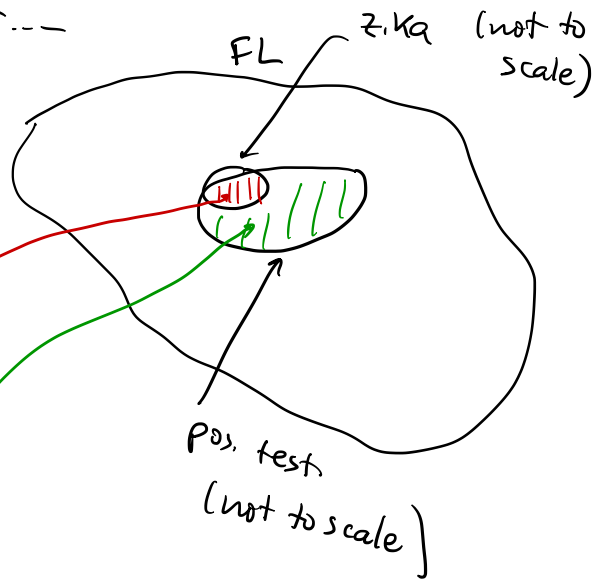
w/ Zika:  $10^7 \cdot 10^{-5} = 10^2 = 100$

w/o Zika: 9,999,900

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test pos w/ Zika:  $100 \cdot 0.99 = 99$

test pos. w/o Zika:  $9,999,900 \times 0.01$   
 $= 99,999$

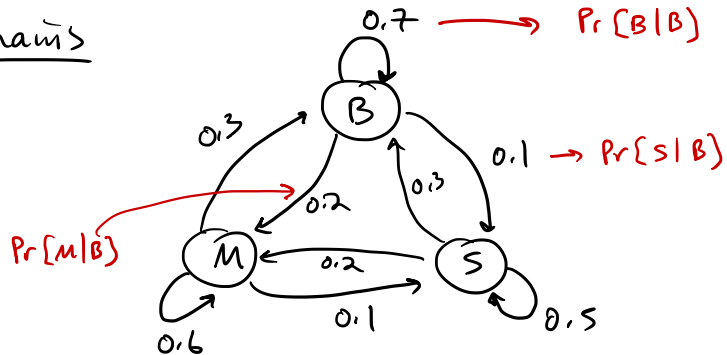


$\therefore$  If you test pos.,

you're about 1,000 times more likely  
to be among 99,999 who don't have Zika  
but test pos. than among the 99  
who do have Zika and test pos.

# Markov chains

B: Bertucci's  
M: Margaritas  
S: Sato



3 reds  $\Rightarrow$  must add  
to 1

