

7 Hiding Implementation Details

Last time we saw the *Specification Design Pattern* as a rather mechanical way to get an implementation that almost out-of-the-box implements a specification for an ADT, using a natural representation via concrete subclasses capturing what each of the creators of the ADT is doing.

That design pattern however reveals a lot of implementation details. The problem with reveal implementation details is that (1) it is information that is not part of the ADT signature, so it shouldn't be available, and (2) if it's there, someone will use it, and if someone uses it, that someone will be in trouble if later on you come back and reimplement the ADT using another implementation that still satisfies the specification but does not reveal quite the same information as your original implementation.

Let's illustrate this with an example, which will come in handy later.

7.1 Another Example: Lists of Integers

Here is a straightforward ADT for lists of integers:

CREATORS

```
empty :      () -> List
cons  :      Int List -> List
```

OPERATIONS

```
isEmpty :    () -> Boolean
first   :    () -> Int
rest    :    () -> List
length :    () -> Int
append  :    List -> List
find    :    Int -> Boolean
isEqual :    List -> Boolean
```

and the obvious specification:

```
empty().isEmpty() = true
cons(n,L).isEmpty() = false
```

`cons(n,L).first() = n`

`cons(n,L).rest() = L`

`empty().length() = 0`

`cons(n,L).length() = 1 + L.length()`

`empty().append(M) = M`

`cons(n,L).append(M) = cons(n,L.append(M))`

`empty().find(f) = false`

$$\text{cons}(n,L).\text{find}(f) = \begin{cases} \text{true} & \text{if } n = f \\ \text{true} & \text{if } L.\text{find}(f) = \text{true} \\ \text{false} & \text{otherwise} \end{cases}$$
$$\text{empty}().\text{isEqual}(M) = \begin{cases} \text{true} & \text{if } M.\text{isEmpty}() = \text{true} \\ \text{false} & \text{otherwise} \end{cases}$$
$$\text{cons}(n,L).\text{isEqual}(M) = \begin{cases} \text{false} & \text{if } M.\text{isEmpty}() = \text{true} \\ \text{true} & \text{if } n = M.\text{first}() \text{ and } L.\text{isEqual}(M.\text{rest}()) \end{cases}$$

Applying the Specification Design Pattern, and providing a reasonable implementation for the canonical methods, we get something like:

```
/*
 * The default result of the Specification Design Pattern
 *
 */

object List {

  def empty () : List = new ListEmpty()
  def cons (n:Int, L:List) : List = new ListCons(n,L)
}

abstract class List {

  def isEmpty () : Boolean
```

```

def first () : Int
def rest () : List
def length () : Int
def append (M:List) : List
def find (f:Int) : Boolean
def isEqual (M:List) : Boolean
}

class ListEmpty () extends List {

  def isEmpty () : Boolean = true

  def first () : Int =
    throw new RuntimeException("empty().first()")

  def rest () : List =
    throw new RuntimeException("empty().rest()")

  def length () : Int = 0

  def append (M:List) : List = M

  def find (f:Int) : Boolean = false

  def isEqual (M:List) : Boolean = M.isEmpty()

  override def equals (other : Any) : Boolean =
    other match {
      case that : List => this.isEqual(that)
      case _ => false
    }

  override def hashCode () : Int = 41

  override def toString () : String = ""
}

class ListCons (n:Int, L:List) extends List {

```

```

def isEmpty () : Boolean = false

def first () : Int = n

def rest () : List = L

def length () : Int = 1 + L.length()

def append (M:List) : List = List.cons(n,L.append(M))

def find (f:Int) : Boolean =
  ((f == n) || L.find(f))

def isEqual (M:List) : Boolean =
  (!(M.isEmpty()) && n==M.first() && L.isEqual(M.rest()))

override def equals (other : Any) : Boolean =
  other match {
    case that : List => this.isEqual(that)
    case _ => false
  }

override def hashCode () : Int =
  41 * (
    41 + n.hashCode()
  ) + L.hashCode()

override def toString () : String = n + " " + L
}

```

Note the structure: a module `List`, an abstract class `List` that *declares* but does not *define* the operations of the ADT,¹ and two concrete classes that are subtypes of the list `List` and that actually define the operations declared in the abstract class (as well as the canonical methods).

Here's some sample testing code to illustrate:

¹A *declaration* is a promise that you will implement the method when you subclass the abstract class. In other words, it's a promise that any subtype of the abstract will supply an implementation for those methods. You do not need to declare the canonical methods, because their canonical: they're always declared by default in every class, abstract or otherwise. A *definition* is a declaration that also supplies some code defining the method.

Now, this way of writing lists is ugly. Let's do something more that uses brackets at the beginning and the end of the list, and commas to separate the elements. It's easy to do if we have a helper method, which I'll call `toStringLst()`, and change `toString()` to use that helper method – in `ListEmpty`:

```
def toStringLst () : String = ""

override def toString () : String = "[]"
```

and in `ListCons`:

```
def toStringLst () : String = ", " + n + L.toStringLst()

override def toString () : String = "[" + n + L.toStringLst() + "]"
```

Of course, if we make this change, then the code does not type-check anymore. Why? Look at the code. In the `toString()` implementation of `ListCons`, we have a call to `L.toStringLst()`, where `L` is a `List`. When the type checker looks at `List` to see if a `toStringLst()` method is declared there, there isn't one. So it gives an error.² Thus, in order to satisfy the type checker, we need to make sure that every subtype of `List` defines `toStringLst()`, which is done by declaring `toStringLst()` in `List`. Here is the resulting implementation of the `LIST` ADT in full

```
/*
 * The default result of the Specification Design Pattern
 * with a nicer toString() method
 *
 */

object List {

  def empty () : List = new ListEmpty()
  def cons (n:Int, L:List) : List = new ListCons(n,L)
}
```

²Recall, the type checker is in charge of ensuring safety — that you never try to call methods that don't exist. Because `List` does not declare `toStringLst()`, someone could presumably add another subtype of `List` that does not define `toStringLst()` (since it's not declared in the abstract class `List`, subtypes of `List` are under no obligation to define it), and that new subtype could be used in the construction of a `List` using `cons`, and when `toString()` on that list tries to call `toStringLst()` on that new subtype of `List`, it will fail to find it, breaking safety.

```

abstract class List {

  def isEmpty () : Boolean
  def first () : Int
  def rest () : List
  def length () : Int
  def append (M:List) : List
  def find (f:Int) : Boolean
  def isEqual (M:List) : Boolean

  def toStringLst () : String
}

class ListEmpty () extends List {

  def isEmpty () : Boolean = true

  def first () : Int =
    throw new RuntimeException("empty().first()")

  def rest () : List =
    throw new RuntimeException("empty().rest()")

  def length () : Int = 0

  def append (M:List) : List = M

  def find (f:Int) : Boolean = false

  def isEqual (M:List) : Boolean = M.isEmpty()

  override def equals (other : Any) : Boolean =
    other match {
      case that : List => this.isEqual(that)
      case _ => false
    }

  override def hashCode () : Int = 41

```

```

def toStringLst () : String = ""

override def toString () : String = "[]"
}

class ListCons (n:Int, L:List) extends List {

  def isEmpty () : Boolean = false

  def first () : Int = n

  def rest () : List = L

  def length () : Int = 1 + L.length()

  def append (M:List) : List = List.cons(n,L.append(M))

  def find (f:Int) : Boolean =
    ((f == n) || L.find(f))

  def isEqual (M:List) : Boolean =
    (!(M.isEmpty()) && n==M.first() && L.isEqual(M.rest()))

  override def equals (other : Any) : Boolean =
    other match {
      case that : List => this.isEqual(that)
      case _ => false
    }

  override def hashCode () : Int =
    41 * (
      41 + n.hashCode()
    ) + L.hashCode()

  def toStringLst () : String = ", " + n + L.toStringLst()

  override def toString () : String = "[" + n + L.toStringLst() + "]"
}

```

We can test this as before:

```

val L1:List = List.cons(33,List.cons(66,List.cons(99,List.empty())))

println("L1 = " + L1)
println("First of L1 = " + L1.first())
println("Rest of L1 = " + L1.rest())
println("Length of L1 = " + L1.length())

```

and we get the arguably nicer output:

```

L1 = [33, 66, 99]
First of L1 = 33
Rest of L1 = [66, 99]
Length of L1 = 3

```

Now, the problem with the above implementation is that it reveals two things: first, that there is a helper function called `toStringLst()`, even though it's not actually part of the ADT, and second, that the implementation is in terms of two concrete classes `ListEmpty` and `ListCons`.

To wit, continuing with the above example:

```

/* accessing the representation directly */
val L2 : List = new List.ListCons(4,L1)
  println("L2 = " + L2)

/* calling a function not in the signature */
println("helper result = " + L1.toStringLst())

```

and we get:

```

L2 = [4, 33, 66, 99]
helper result = , 33, 66, 99

```

That's bad: if later on we decide on a new implementation of `toString()` and get ride of `toStringLst()`, then anyone that relied on that helper function existing will have their code suddenly stop working. The function is not in the signature, so it should not be available to anyone else. Similarly, if later on we decide on a new implementation of lists altogether, one that does not rely on two concrete classes, but perhaps on three, or just one, not necessarily named `ListCons`, then anyone creating a direct form of `ListCons` will have their code suddenly stop working.

So: how do we hide the helper method, and how do we hide the fact that there are those two representation classes?

Let's deal with the second question first, partly because we cannot really hide helper functions in an abstract/concrete class combo without doing this first.

7.2 Hiding Representation Classes

So how do we hide information. The main tool we have for hiding information is to make that information private to some area of the code. For instance, if we mark a method as `private`, as in:

```
class Foo1 {  
  
    def method1 (x:Int):Int = x+method2(x)  
  
    private def method2 (x:Int):Int = x*2  
}
```

then `method2` is only visible from within an instance of `Foo1`, while `method1` is visible from outside as well as from within. Thus, while this works:

```
val f = new Foo1  
f.method1(10)
```

this causes a compile-time error saying that `method2` is inaccessible:

```
val f = new Foo1  
f.method2(10)
```

We can similarly make fields private. For instance,

```
class Bar (x:Int) {  
  
    val field1:Int = x+1  
  
    private val field2:Int = 2*x  
}
```

and the following works:

```
val b = new Bar(5)  
b.field1
```

while the following fails for the same reason as `f.method2(10)` failed earlier:

```
val b = new Bar(5)  
b.field2
```

So in general, we can make any component of a class private and thus hide it from whatever is outside the class. Thus, to hide the representation classes `ListEmpty` and `ListCons`, then, it turns out we can simply them component of `List`! Such classes, nested inside other classes, are called *nested classes*.³ Here is a simple example of a nested class, not hidden:

```
object Foo2 {  
  
  def method1 (x:Int):Bar = new Bar(2*x)  
  
  class Bar (value:Int) {  
  
    def method2 (y:Int):Int = value+y  
  }  
}
```

If we call `method1` in `Foo2`, the result will be an instance of class `Foo2.Bar`:

```
val b = Foo2.method1(10)  
b.method2(30)
```

and the result should be 50. Note the type of `b`: `Foo2.Bar` — it has type `Bar` defined inside of `Foo2`. Classes accessed just like fields or methods when they occur inside `Foo2`. We can create instances of `Foo2.Bar` directly too:

```
val b = new Foo2.Bar(10)  
b.method2(30)
```

which returns value 40.

Now, if we hide the nested class `Bar`:

```
object Foo3 {  
  
  def method1 (x:Int):Bar = new Bar(2*x)  
  
  private class Bar (value:Int) {  
  
    def method2 (y:Int):Int = value+y  
  }  
}
```

³Following Java-based terminology, a class defined inside a module is sometimes called a *nested class*, while a class defined inside another class is a special kind of nested class called an *inner class*. Inner classes are exceedingly expressive, and are related to closures. In other words, inner classes give you `lambda`.

then we get a problem trying to compile because our hiding worked too well: we do not have access to type `Bar` anymore (since `Bar` is hidden) and the system does not know how to refer to the value returned by `method1()` — the Scala compiler complains that the private class `Bar` *escapes* the class in which it is defined to be private. So one trick there is to simply use an abstract class to tell the compiler that yes, there is such a class, and it implements the following functions, but the implementation remains hidden:

```
object Foo4 {

  def method1 (x:Int):Bar = new BarImplementation(2*x)

  private class BarImplementation (value:Int) extends Bar {

    def method2 (y:Int):Int = value+y
  }
}

abstract class Bar {

  def method2 (y:Int):Int
}
```

Now we can still create instances of `BarImplementation` by calling `Foo4.method1()`, but not directly by calling `new Foo4.Bar()`.

Make sure you understand how the above works, since it's the structure we are going to be using. *Question: what happens if you make `FOOx` above a class instead of an object? Try it — see how you can make it work. Again, use the analogy that such defined classes are accessed just like fields or methods.*

Here is the above structure applied to our implementation of the LIST ADT — where I've compressed the nested classes to make the structure more apparent.

```
/*
 * The default result of the Specification Design Pattern
 *   with a nicer toString() method
 *   modified to hide concrete representation classes
 *
 */

object List {

  def empty () : List = new ListEmpty()
```

```

def cons (n:Int, L:List) : List = new ListCons(n,L)

private class ListEmpty () extends List {
  def isEmpty () : Boolean = true
  def first () : Int =
    throw new RuntimeException("empty().first()")
  def rest () : List =
    throw new RuntimeException("empty().rest()")
  def length () : Int = 0
  def append (M:List) : List = M
  def find (f:Int) : Boolean = false
  def isEqual (M:List) : Boolean = M.isEmpty()

  override def equals (other : Any) : Boolean =
    other match {
      case that : List => this.isEqual(that)
      case _ => false
    }
  override def hashCode () : Int = 41
  override def toString () : String = "[]"
  def toStringLst () : String = ""
}

private class ListCons (n:Int, L:List) extends List {
  def isEmpty () : Boolean = false
  def first () : Int = n
  def rest () : List = L
  def length () : Int = 1 + L.length()
  def append (M:List) : List = List.cons(n,L.append(M))
  def find (f:Int) : Boolean =
    ((f == n) || L.find(f))
  def isEqual (M:List) : Boolean =
    (!(M.isEmpty()) && n==M.first() && L.isEqual(M.rest()))

  override def equals (other : Any) : Boolean =
    other match {
      case that : List => this.isEqual(that)
      case _ => false
    }
  override def hashCode () : Int =
    41 * (

```

```

        41 + n.hashCode()
    ) + L.hashCode()

    override def toString () : String = "[" + n + L.toStringLst() + "]"
    def toStringLst () : String = ", " + n + L.toStringLst()
  }
}

abstract class List {

    def isEmpty () : Boolean
    def first () : Int
    def rest () : List
    def length () : Int
    def append (M:List) : List
    def find (f:Int) : Boolean
    def isEqual (M:List) : Boolean

    def toStringLst () : String
}

```

7.3 Hiding Helper Functions

Now there only remain the question of how to hide the helper function.

If a helper function was only needed in one particular class, then hiding that helper function is easy — it's just like hiding `method2()` in `Foo1` above: `method2()` is only needed inside `Foo1`, by method `method1()` in particular, so it is marked `private`. So hiding a helper function used only within a class is trivial.

The problem, as we saw, is that `toStringLst()` is called across the two concrete classes that are subtypes of `List`, and because of that, we had to declare it in the abstract class `List`. So we want to say that `toStringLst()` is only available within `List` and its subtypes. The key thing here is *and its subtypes*. The annotation we want here is `protected`, which is just like `private`, except it also makes a declaration visible to the subtypes. (Because declarations in an abstract class *have* to be visible from the subtypes (that's the whole point of declarations in abstract classes), the Scala compiler will not allow you to mark a declaration as `private` — you have to use `protected`.) And if we do that, it works:

```

/*
 * The default result of the Specification Design Pattern

```

```

*   with a nicer toString() method
*   modified to hide concrete representation classes
*   and to hide helper functions
*
*/

object List {

  def empty () : List = new ListEmpty()
  def cons (n:Int, L:List) : List = new ListCons(n,L)

  private class ListEmpty () extends List {
    def isEmpty () : Boolean = true
    def first () : Int =
      throw new RuntimeException("empty().first()")
    def rest () : List =
      throw new RuntimeException("empty().rest()")
    def length () : Int = 0
    def append (M:List) : List = M
    def find (f:Int) : Boolean = false
    def isEqual (M:List) : Boolean = M.isEmpty()

    override def equals (other : Any) : Boolean =
      other match {
        case that : List => this.isEqual(that)
        case _ => false
      }
    override def hashCode () : Int = 41
    override def toString () : String = "[]"
    def toStringLst () : String = ""
  }

  private class ListCons (n:Int, L:List) extends List {
    def isEmpty () : Boolean = false
    def first () : Int = n
    def rest () : List = L
    def length () : Int = 1 + L.length()
    def append (M:List) : List = List.cons(n,L.append(M))
    def find (f:Int) : Boolean =
      ((f == n) || L.find(f))
    def isEqual (M:List) : Boolean =

```

```

        (!(M.isEmpty()) && n==M.first() && L.isEqual(M.rest()))

override def equals (other : Any) : Boolean =
    other match {
        case that : List => this.isEqual(that)
        case _ => false
    }
override def hashCode () : Int =
    41 * (
        41 + n.hashCode()
    ) + L.hashCode()
override def toString () : String = "[" + n + L.toStringLst() + "]"
def toStringLst () : String = ", " + n + L.toStringLst()
}
}

abstract class List {

    def isEmpty () : Boolean
    def first () : Int
    def rest () : List
    def length () : Int
    def append (M:List) : List
    def find (f:Int) : Boolean
    def isEqual (M:List) : Boolean

    // protected to make it unavailable outside List and its subtypes
    protected def toStringLst () : String
}

```

And there: we cannot create instances of `ListEmpty` and `ListCons` from outside `List` (and therefore we have to use the creators to create lists), and once we have a list we cannot call its `toStringLst` function, as the compiler will fail. For instance, if we try to compile and run:

```

val L1:List = List.cons(33,List.cons(66,List.cons(99,List.empty())))

println("L1 = " + L1)
println("First of L1 = " + L1.first())
println("Rest of L1 = " + L1.rest())

```

```
println("Length of L1 = " + L1.length())

/* calling a function not in the signature */
println("helper result = " + L1.toStringLst())
```

Then the compiler complains

```
<console>:26: error: method toStringLst cannot be accessed in List
  println('helper result = ' + L1.toStringLst())
                                ^
one error found
```

As desired.

Now, note something interesting. The above *should not* work! We're declaring `toStringLst()` protected in `List`. Which is the same essentially as making it `private` — it means, in particular, that there we are not allowed to refer to that method from outside the class `List` itself. But we're invoking `toStringLst()` on something of type `List` in `ListCons` in method `toString()`. Why does the compiler let us do that? The answer is that we're made use of a particular back-door in Scala. The fact that module in which `ListEmpty` and `ListCons` are defined has the same name `List` as the abstract class is the key point — as I mentioned earlier in the course, they are called *companions*: a module and a class of the same name. Being companions means that they can access each other's private components as though they were defined within themselves. So in the case of `List`, for the purposes of what's private/public, it's as if we had defined:

```
object/class List {

  def empty () : List = new ListEmpty()
  def cons (n:Int, L>List) : List = new ListCons(n,L)

  private class ListEmpty () extends List {
    ...
  }

  private class ListCons (n:Int, L>List) extends List {
    ...
  }

  def isEmpty () : Boolean
  def first () : Int
  def rest () : List
  def length () : Int
```



```
def append (M:List) : List
def find (f:Int) : Boolean
def isEqual (M:List) : Boolean

// protected to make it unavailable outside List and its subtypes
protected def toStringLst () : String
}
```

meaning that whatever is inside this combo `List` can access `toStringLst()` since it is defined within the same scope. And since `ListCons` is inside the scope (it is a component within `List`), it can access `toStringLst()`.

This works only because the module and the class have the same name `List`. If we change that, that is, if we change the name of module `List` to something like `ListCreators`, this breaks miserably: the system complains that `toString()` in `ListCreators.ListCons` attempts to call the inaccessible method `toStringLst()` in `List`. If you ever get that error, you'll know where to look.