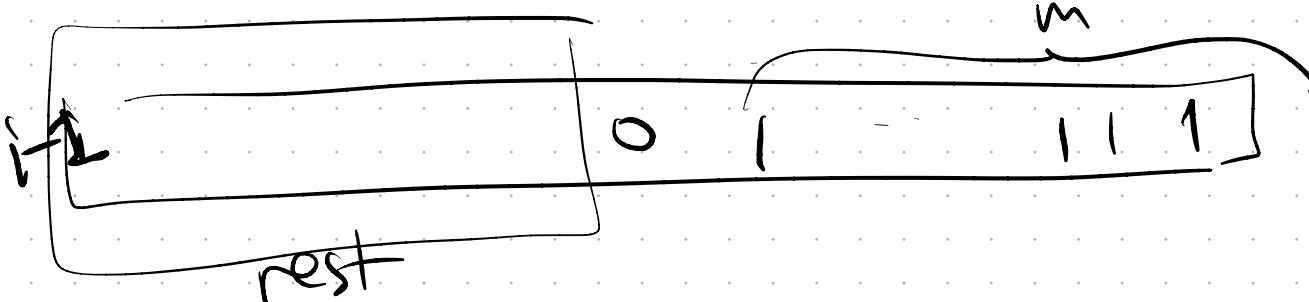
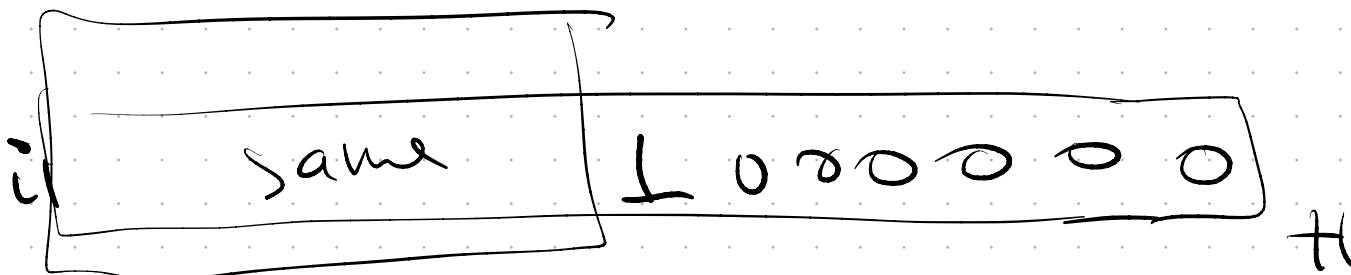


$$fm \cos t = mH$$

$$\hat{C}_i = C_i + \Delta\phi_i = mH + (1-m) = 2$$

$$\Delta\phi = 1 - m$$



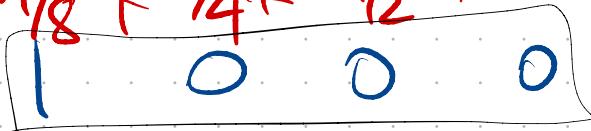
Avgortized Analysis

Binary counter 6 bits operation $i+1$ cost = # bits changed

$b_5 \ b_4 \ b_3 \ b_2 \ b_1 \ b_0$

global # changes in n ops

$$\frac{n}{10} + \frac{n}{8} + \frac{n}{4} + \frac{n}{2} + n \leq 2n$$



$\delta \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0$

$\delta \quad 0 \quad 0 \quad 0 \quad 1 \quad 0 \quad 1$

$\delta \quad 0 \quad 0 \quad 0 \quad 1 \quad 0 \quad 0$

$\delta \quad 0 \quad 0 \quad 0 \quad 0 \quad 1 \quad 1$

$\delta \quad 0 \quad 0 \quad 0 \quad 0 \quad 1 \quad 0$

$\delta \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 1$

$\delta \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0$

true cost

global count method

$$\phi(T_i) =$$

= # of "1" bits

$$1 \quad \hat{c}_i = c_i + \Delta$$

3 2

2 2

2 2

1 2

2 2

1 2

4

1

2

1

3

1

2

1

0

worst case for op

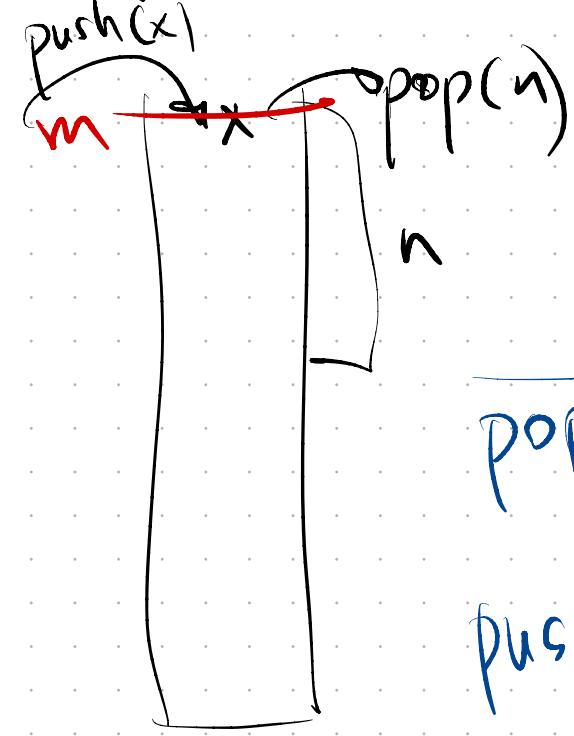
$$n \text{ operations} \times \Theta(\log n) = \Theta(n \log n)$$

Naive analysis: n operations $\times \Theta(\log n) = \Theta(n \log n)$

Stack operations

$\text{push}(x) \rightarrow \text{push } x \text{ to top}$

$\text{pop}(n) \rightarrow \text{pops } n \text{ value from top}$
(or all if $n \leq \text{avail}$)



n operations

- valid stack

naive runtime

$$n \times \text{pop}(n) = \Theta(n^2)$$

• accounting method

	true cost c_i	amortized cost \hat{c}_i (savings)
$\text{POP}(n)$	n	0
push	1	2

c_i = true cost of op_i

\hat{c}_i = amortized cost of op_i

K-ops

Guarantees:

$$\sum_{i=1}^k c_i \leq \sum_{i=1}^k \hat{c}_i$$

What we payed so far (k) is enough to cover actual exp.

$$\phi(\text{stack}) = \text{size of stack} = m \quad \Delta$$

$$\overbrace{\text{push}}^{\Delta\phi} = \text{push} + \Delta\phi_{\text{push}} = 1 + 1 = 2$$

$$\overbrace{\text{pop}(k)}^{\Delta\phi_{\text{pop}}} = \text{pop}(k) + \Delta\phi_{\text{pop}} = k + (-k) \Rightarrow \xrightarrow{k \text{ lower}}$$

• Potential (situation
danger
datastructure)
!

T_i = datastructure / situation
after i -th operation

$$\phi(T_i) \geq 0 \quad \phi(T_0) = 0$$

Th $\Sigma_{i=1}^k c_i$ happens $\forall i \Rightarrow$ Guarantee works.

Proof: $\sum_{i=1}^k \hat{c}_i \geq \sum_{i=1}^k c_i + \phi(T_i) - \phi(T_{i-1}) = \sum_{i=1}^k c_i +$

$$\cancel{\phi(T_k) - \phi(T_{k-1})} +$$

$$\cancel{\phi(T_{k-1}) - \phi(T_{k-2})} +$$

$$= \sum_{i=1}^k c_i + \cancel{\phi(T_k) - \phi(T_0)}$$

suf cond $\hat{c}_i - c_i \geq \Delta \phi_i$

$$\hat{c}_i \geq c_i + \phi(T_i) - \phi(T_{i-1})$$

$\Delta \phi_i$ = change in
potential due to
 ϕ_i