

SCG Example Labs

Ahmed Abdelmeged

Karl Lieberherr

Structures of SCG

- These examples are to be read with the SCG paper as background.
- The best way to represent Domain, Lab and Claim is to have Domain and Lab as top-level classes and Claim nested inside Lab. Lab has a Domain as field to give all claims access to Domain functionality.
- Instance and Solution are nested inside Domain.

Structures of SCG

- We use a Java-like syntax but the goal is to use only one or two lines per item for those simple introductory labs.

Structures of SCG

```
Domain  
Instance  
Solution  
valid(i: Instance, s: Solution)  
quality(i: Instance, s: Solution)
```

```
Lab  
d : Domain  
proto: Protocol  
Claim  
  claim parameters  
isp(i:d.Instance)  
p(I:d.Instance[],S:d.Solution[])  
stronger(c2: Claim)  
distance(c2: Claim)
```

Calculus Lab

- Have your students mastered calculus (minimizing and maximizing functions)?
- The next lab shows a lab to test their skills.

Structures of SCG

NEW

Domain

Instance

Solution

valid(i: Instance, s: Solution)

quality(i: Instance, s: Solution)

Example: calculus problem

SaddlePoint

Instance = [0,1]

Solution = [0,1]

valid(i,s) = true

quality(i,s) = $i*s + (1-i)*(1-s^2)$

Lab

d: Domain

proto: Protocol

Claim

claim parameters

isp(i:d.Instance)

p(I:d.Instance[],S:d.Solution[])

stronger(c2: Claim)

distance(c2: Claim)

SaddlePointLab

SaddlePoint

O:I[0], P:S[1] of I[0]

SaddlePointLabClaim

q: [0,1]

isp(i)=true

p(I,S)=d.quality(I[0],S[1])>=q

stronger(c2) = q>c2.q

distance(c2) = q-c2.q

new SaddlePointLab.Claim(q=0.6)

Programming an Algorithm

- Have your students understood the Gale-Shapley algorithm?
- Next come two labs where they can demonstrate their skills through their avatar in a full-round-robin tournament.
 - In the first lab they test each other's programs to see whether they match each other's best solution.
 - In the second lab, they find worst-case inputs.

Structures of SCG

Domain
Instance
Solution
valid(i: Instance, s: Solution)
quality(i: Instance, s: Solution)

Example: GS algorithm

GaleShapleyBasic
Instance = Preferences
Solution = Assignment
valid(i, s) = s is syntactically correct for i
quality(i, s) = s is semantically correct for i
1 if true, 0 if false.

Lab
d: Domain
proto: Protocol
Claim
claim parameter definitions
isp(i:d.Instance)
p(I:d.Instance[],S:d.Solution[])
stronger(c2: Claim)
distance(c2: Claim)

GaleShapleyBasicLab
GaleShapleyBasic
O:I[0], P:S[1] of I[0], O:S[2] of I[0]
GSAtLeastAsGoodAsYouClaim
none
isp(i)=true
 $p(I, S) = d.\text{quality}(I[0], S[1]) \geq d.\text{quality}(I[0], S[2])$
stronger(c_2) = false
distance(C_2) = 0

NEW GSAtLeastAsGoodAsYouClaim()

Structures of SCG

```
Domain  
Instance  
Solution  
valid(i: Instance, s: Solution)  
quality(i: Instance, s: Solution)
```

Example: Worst-Case of GS algorithm

```
GaleShapley (GS)  
Instance = Nat //number of people  
Solution = Preferences  
valid(i,s) = s is syntactically correct for i  
quality(i,s) = GS iterations for s and i
```

```
Lab  
d: Domain  
proto: Protocol  
Claim  
claim parameters  
isp(i:d.Instance)  
p(I:d.Instance[],S:d.Solution[])  
stronger(c2: Claim)  
distance(c2: Claim)
```

```
GaleShapleyWorstCaseLab  
GaleShapley  
O:I[0], P:S[1] of I[0]  
GSWCLClaim  
n:Nat, q:Nat  
isp(i)=(i=n) //singleton  
p(I,S)=d.quality(I[0],S[1])>= q  
stronger(c2)=this.q>c2.q  
distance(c2)=this.q-c2.q
```

```
new GSWCLClaim(n=10,q=30)
```

Maximum Satisfiability

- The next lab is about a paper by David Johnson in the 1970's which is covered now in some algorithm text books, like Kleinberg and Tardos.
- The following MaxSat lab covers several skills, such as working with randomized algorithms and then derandomizing them.

Structures of SCG

Domain

Instance

Solution

valid(i: Instance, s: Solution)

quality(i: Instance, s: Solution)

Example: MaxSat

Satisfiability

CNF

Assignment

$\text{valid}(i,s) = \text{all variables in } i \text{ assigned once}$

$\text{quality}(i,s) = \text{fraction of satisfied clauses in } i$

Lab

d: Domain

proto: Protocol

Claim

claim parameters

isp(i:d.Instance)

p(I:d.Instance[], S:d.Solution[])

stronger(c2: Claim)

distance(c2: Claim)

MaxSatLab

Satisfiability

$O:I[0], P:S[1] \text{ of } S[0]$

MSLClaim

$q:[0,1], k:\text{Nat} \text{ (clause length)}$

$\text{isp}(i)=\text{clauses in } i \text{ have length } \geq k$

$p(I,S)=d.\text{quality}(I[0],S[1]) \geq q$

$\text{stronger}(c2)=q > c2.q$

$\text{distance}(c2)=q - c2.q$

new MSLClaim(q=1-(1/2³), k=3)

Generalized MaxSat

- The next lab is based on a paper by Lieberherr and Specker (2012).

Structures of SCG

Domain

Instance

Solution

valid(i: Instance, s: Solution)

quality(i: Instance, s: Solution)

Example: Boolean GeneralizedMaxSat = BMaxCSP

BooleanCSP

Sequence of Boolean constraints

Assignment

$\text{valid}(i,s) = \text{all variables in } i \text{ assigned once}$

$\text{quality}(i,s) = \text{fraction of sat. constraints in } i$

Lab

d: Domain

proto: Protocol

Claim

claim parameters

isp(i:d.Instance)

p(I:d.Instance[],S:d.Solution[])

stronger(c2: Claim)

distance(c2: Claim)

BooleanMaxCSPLab

BooleanCSP

$O:I[0], P:S[1] \text{ of } I[0]$

GenBooleanMaxSatClaim

$q:[0,1], r:\{R1,R2,\dots\}$

$\text{isp}(i)=\text{constraints in } i \text{ use only } r$

$p(I,S)=d.\text{quality}(I[0],S[1])>=q$

$\text{stronger}(c2)=q>c2.q$

$\text{distance}(c2)=q-c2.q$

`new GenBooleanMaxSatClaim(q=0.618, r={R1,R2})`

Local to Global

- The next lab is based on several papers inspired by a JACM paper by Lieberherr and Specker in 1981.
- The lab is about studying how local properties of a conjunctive normal form translate into global properties.

Structures of SCG

Domain

Instance

Solution

valid(i: Instance, s: Solution)

quality(i: Instance, s: Solution)

Example: BooleanMaxCSPLocalGlobal

BooleanCSP

Sequence of Boolean constraints

Assignment

$\text{valid}(i,s) = \text{all variables in } i \text{ assigned once}$

$\text{quality}(i,s) = \text{fraction of sat. constraints in } i$

Lab

d: Domain

proto: Protocol

Claim

claim parameters

isp(i:d.Instance)

p(I:d.Instance[], S:d.Solution[])

stronger(c2: Claim)

distance(c2: Claim)

BooleanMaxCSPLab

BooleanCSP

O:I[0], P:S[1] of I[0]

BMCLClaim

q:[0,1], r:{R1,R2,...}, k:Nat

isp(i)=(constraints in i use only r)

and (any k constraints in i are satisfiable)

p(I,S)=d.quality(i[0],s[1])>=q

stronger(c2)=q>c2.q

distance(c2)=q-c2.q

new BMCLClaim(q=0.618, r={R1,R2,R3,R4}, k=2)

Manufacturing Lab

- The next lab is about an efficient manufacturing problem where raw materials are turned into a product.
- The lab is underspecified in that the details about the `isp` function are missing.

Structures of SCG

Domain
Instance
Solution
valid(i: Instance, s: Solution)
quality(i: Instance, s: Solution)

Example: Solar Cells

SolarCells
RawMaterials
Product
 $\text{valid}(i,s) = \text{only raw materials used}$
 $\text{quality}(i,s) = \text{energy efficiency of } s \text{ for } i$

Lab
d: Domain
proto: Protocol
Claim
claim parameters
isp(i:d.Instance)
p(I:d.Instance[],S:d.Solution[])
stronger(c2: Claim)
distance(c2: Claim)

SollarCellsLab
SollarCells
 $O:I[0], P:S[1] \text{ of } I[0]$
SCLClaim
 $q:[0,1], k:\text{Nat} \text{ (raw material parameter)}$
 $\text{isp}(i)= \dots$
 $p(I,S)=d.\text{quality}(I[0],S[1])> q$
 $\text{stronger}(c2)=q>c2.q$
 $\text{distance}(c2)=q-c2.q$

new SCLClaim(q=0.7,k=3)

Lab Reductions

- With the next example we show the usefulness of lab reductions. A lab L1 reduces to a lab L2 ($L1 < L2$) if a defense strategy for the claims in L2 guarantees a defense strategy for the claims in L1. Ideally, the claims in L2 are simpler.
- L1 reduces to L2 if we can use a black box for L2 to solve L1. The black box makes all perfect decisions, including claims it can defend.

Lab Reductions

- A mapping from L1 to L2 is a computable function $f : \text{Domain} \rightarrow \text{Claim}$ such that for any
 - L1.Domain \rightarrow L2.Domain
 - L1.Claim \rightarrow L2.Claim
 - propose
 - oppose/agree
 - provideInstance
 - solveInstance
 - refute

expression in a, d, n using multiplication, addition and division.

To simplify, replace $(1+2+\dots+n)$ by $n*(n+1)/2$.

$$\sum_{k=1}^n a + dk = (a + d) + (a + 2d) + (a + 3d) + \dots + (a + nd) = na + (1 + 2 + 3 + \dots + n)d$$

Structures of SCG

Domain with name

Instance

Solution

valid(i: Instance, s: Solution)

quality(i: Instance, s: Solution)

Lab with name

d: Domain

claim parameter definitions

instance set predicate

refutation predicate

protocol

stronger(c1,c2: Claim)

distance(c1,c2: Claim)

Example: Arithmetic Sequences Sum

ArithmeticSequences2

triple a,d,n: Nat

expression in a,d,n

valid(i,s) = s uses +,*,/ and vars in i

quality(i,s) = 1 if s is correct for i

Claim with name

lab : Lab

claim parameter values

Structures of SCG

Domain with name

Instance

Solution

valid(i: Instance, s: Solution)

quality(i: Instance, s: Solution)

Lab with name

d: Domain

claim parameter definitions

instance set predicate

refutation predicate

protocol

stronger(c1,c2: Claim)

distance(c1,c2: Claim)

Claim with name

lab : Lab

claim parameter values

Example: Arithmetic Sequences Sum

ArithmeticSequences

expression in a,d,n: Nat uses +,*,/

assignment to a,d,n

valid(i,s) = s assigns all 3 variables

quality(i,s) = 1 iff i gives correct sum for s

ArithmeticSequencesLab

ArithmeticSequences

none

singleton

quality(i[0],s[1])=1 true

O:i[0], P:s[1] of s[0]

false

0

ASLClaim(

ArithmeticSequencesLab,

) //none

Structures of SCG

Domain with name

Instance

Solution

valid(i: Instance, s: Solution)

quality(i: Instance, s: Solution)

Lab with name

d: Domain

claim parameter definitions

instance set predicate

refutation predicate

protocol

stronger(c1,c2: Claim)

distance(c1,c2: Claim)

Claim with name

lab : Lab

claim parameter values

Example: Arithmetic Sequences Sum

ArithmeticSequencesInduction

sum[k=1..n] 2+3k = 2n+3(n(n+1))/2

sequence of steps: induction proof

valid(i,s) = proof s is syntactically correct

quality(i,s) = 1 iff proof is correct

ArithmeticSequencesInductionLab

ArithmeticSequencesInduction

equation

singleton

quality(i[0],s[1])=1

O:i[0], P:s[1] of s[0]

false

0

ASLClaim3(

Alice,

ArithmeticSequencesInductionLab,

sum[k=1..n] 2+3k = 2n+3(n(n+1))/2

Structures of SCG

Domain with name

Instance

Solution

valid(i: Instance, s: Solution)

quality(i: Instance, s: Solution)

Example: HighestSafeRung

HighestSafeRung

pair(n,k)

decision tree

valid(i,s) = s is correct for (n,k)

quality(i,s) = depth(s)

Lab with name

d: Domain

claim parameter definitions

instance set predicate

refutation predicate

protocol

stronger(c1,c2: Claim)

distance(c1,c2: Claim)

HighestSafeRungLab

HighestSafeRung

n->Nat,k->Nat,q->Nat

singleton {(n,k)}

quality(i[0],s[1])<=q

O:i[0], P:s[1] of s[0]

c1.q<c2.q

c1.q-c2.q

Claim with name

lab : Lab

claim parameter values

HSRClaim(

HighestSafeRungLab,

n->25,k->2,k->5)

lower quality
is better

Structures of SCG

Domain with name

Instance

Solution

valid(i: Instance, s: Solution)

quality(i: Instance, s: Solution)

Lab with name

d: Domain

claim parameter definitions

instance set predicate

refutation predicate

protocol

stronger(c1,c2: Claim)

distance(c1,c2: Claim)

Example: LeafCovering

LeafCovering

Set of trees. Set M=subset of GCP of trees.

witness (leaf in GCP) of non-coverage by M

valid(i,s) = s is correct for i

quality(i,s) = unused

Claim with name

proponent: Scholar

lab : Lab

claim parameter values

Structures of SCG

Domain with name

Instance

Solution

valid(i: Instance, s: Solution)

quality(i: Instance, s: Solution)

Lab with name

d: Domain

claim parameter definitions

instance set predicate

refutation predicate

protocol

stronger(c1,c2: Claim)

distance(c1,c2: Claim)

Claim with name

proponent: Scholar

lab : Lab

claim parameter values

Example: LeafCovering

LeafCovering

LeafCoveringProblem : Set of trees. Set M=subs

Program

valid(i,s) = s is correct for i

quality(i,s) = unused

LeafCoveringLab

LeafCovering

m: Nat (size of M)

instanceSetP(i,m)= |i.M|=m

quality(i[0],s[1])<=q

O:i[0], P:s[1] of s[0]

c1.q>c2.q

c1.q-c2.q

HSRClaim(

Alice,

HighestSafeRungLab,

25,2,5)

SCG Truth Table Interpretation

- no competition: P 0 and O 0 everywhere
- not fair: the player different from the “not perfect” player loses a point

claim	dec	out	P	O	cB	aB	oB	Blame Justification
F *	a	sO	0	0	P	O	-	O did not refute a claim it disputed
T	a	sO	0	0	-	-	-	
F *	d	sP	1	-1	P	-	O	O failed to support a claim it agreed with
T	d	sP	1	-1	-	O	O	
F	a	rP	1	-1	P	O	O	P failed to support a claim it proposed
T *	a	rP	1	-1	-	-	O	
F	d	rO	-1	1	P	-	P	P failed to support a claim it proposed
T *	d	rO	-1	1	-	O	P	

claim	dec	out	P	O	cB	aB	oB	Blame Justification
F *	a	sO	<i>paso</i>	<i>oaso</i>	P	O	-	O did not refute a claim it disputed
T	a	sO			-	-	-	
F *	d	sP	<i>pdsp</i>	<i>odsp</i>	P	-	O	O failed to support a claim it agreed with
T	d	sP			-	O	O	
F	a	rP	<i>parp</i>	<i>oarp</i>	P	O	O	P failed to support a claim it proposed
T *	a	rP			-	-	O	
F	d	rO	<i>pdro</i>	<i>odro</i>	P	-	P	P failed to support a claim it proposed
T *	d	rO			-	O	P	